

Equivalent answers: (smooth) "irreducible non-sinsular projective alsobraic curve(overk) of genus 1 furnished with a point O"

This might still not be explicit, but if  $char(K) \neq 2,3$  (as in mort cases) we have the following explicit definition:

Def 1 An elliptic curve over a field K (with char(K)  $\notin$  {2,33) is a plane algebraic curve given by E:  $y^2 = x^3 + Ax + B$ With A, B  $\in$  K and  $4A^3 + 27B^2 \pm 0$ . Example: E:  $y^2 = x^3 - 25x$  (9.6) E(K) real pt. E(R) 5x (9.6) K-rational pts' rational points  $E(Q) \ni (-9,6)$ 

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 Elliptic corves twined out to be useful to answer classical mathematical questions. Nort famous example: Fermat's last theorem: For nz3 at b = c has no integer solutions a,b,c ∈Z with a.b.c ≠ 0.
 Frey ~> If there is a solution the the elliptic curve after change of variables = c (x-a)(x+b)

Tauiyama Conj: Every E/Q is modular. - Shimura Conj: Every E/Q is modular. Wiles (1994): This is true! => FLT (+Taylor) "Modularity theorem"  Elliptic curves also have proclical applications in Cryptography, factoring integers, etc... (later).

We start by talking about general diophantine equations: (named after Diophantus of Alexandria)  $C: f(X_1, X_2, \dots, X_r) = O$  $f(x_1, x_r) \in \mathbb{Z}[x_1, \dots, x_r]$ ,  $\deg(f) = n$ Natural guestions: a) Are the rational or integer solutions? b) If so, can we find them? c) If we have solutions can we find more? d) Can we find all? ~> Determine C(Z), C(Q).

Case r=1 variable (any degree)  $C: f(x) = a_n x'' + \dots + a_i x_i + a_j \quad a_i \in \mathbb{Z}$ Lemma:  $|f x = f \in Q, f(x) = \sigma$  $\Rightarrow p | a_0, q | a_n.$ ~> For giren ao,..., an one just needs to check finitely many X. Case r=2 variables, derree n=1  $C: a_X + b_Y = C , a_b \neq 0$ · Infinitely many solutions over Q (For any  $X \in \mathbb{Q}$ ,  $Y = \frac{C - \alpha X}{b}$ )

. Over Z: Solution over Z (=) gcd (ab) (C (Euklidr alsorithm/Bezout's lemma)

Case V=2 var, desree n=2 (Conicr)  
C: 
$$ax^2+bxy+cy^2+dx+ey+f=0$$
  
(For a):  
Over Q: "local-to-global" principle:  
Thm (Hasse-Minkowki Theorem)  
C has a Q-pt  $\iff$  C has points "locally"  
(C(Q)  $\neq \emptyset$ ) at all "places":  
C(R)  $\neq \emptyset$  and  
(See [L] Appendix C) p-adis C(Qp)  $\neq \emptyset$  for  
all P.  
"solutions mod p<sup>n</sup> for all n".  
Example: C:  $x^2+y^2+1=0$ , C(R)=C(Q)= $\emptyset$   
• C:  $x^2+y^2-3=0$  has  
No vational solution since it has  
No solution mod 4

Here 
$$C(Q) = \emptyset = C(Q_{4})$$
  
(but  $C(R) \neq \emptyset$ )  
 $X^{2} + y^{2} = 113$  has solution  $(x_{1}y_{1}) = (7,8)$ .  
Can we find more?  
Line  
 $Valional$   
 $Valio$ 

<u>Case r=2 var, des f n=3</u> (plane cubic)  $C: ax^{2} + bx^{2} + Cxy^{2} + \dots + j = 0$ In n=2 care we had either no Q-pt or ∞-many. For n=3 we can also just have finitely many Q-pts. But everything is much harder! • the local - to - global principle doesn't work: There are cubics with rol's over IR, Qp bot no Q-rol. "Selmer's example"  $3x^3 + 4y^3 = 5$ . No algorithm to find all Q-pts (even if we have one) laler (projective) -Elliptic curve: non-ringular plane cubic curve with at least one Q-pt. ~> can always brins it in the form y'= X+Ax+B